Amendments to the Claims:

This listing of claims will replace all prior versions, and listings, of claims in the application.

Listing of Claims:

1. (Currently Amended) A method for recovering data in a storage device comprising:

providing a first base block copy having a first sequence number, a first copy pointer and a first checksum;

providing a second base block copy having a second sequence number, a second copy pointer and a second checksum;

defining in the storage device information related to a first data structure within a plurality of copies of a second data structure, wherein the plurality of copies includes the first base block copy and the second base block copy; and

rebuilding the information related to the first data structure using the first base block copy ~~a copy of one of the plurality of copies of the second data structure~~ upon corruption of the second base block copy ~~another copy of one of the plurality of copies of the second data structure~~.

2. (Original) The method as claimed in claim 1 wherein: the storage device performs a write operation; and defining the information related to a first data structure includes updating the plurality of copies of the second data structure prior to the write operation.

3. (Original) The method as claimed in claim 1 wherein defining the information related to a first data structure includes differentiating which of the plurality of copies of the second data structure has the most recent data.

4. (Original) The method as claimed in claim 1 wherein defining the information related to a first data structure includes selecting each of the plurality of copies of the second data structure so at most one of the plurality of copies of the second data structure can be corrupted.

5. (Previously Presented) The method as claimed in claim 1 wherein: the first data structure has an intended state; and rebuilding the information related to the first data structure includes using one or more of the plurality of copies of the second data structure to rebuild the intended state of the first data structure after the corruption of one of the plurality of copies of the second data structure.

6. (Previously Presented) The method as claimed in claim 1 wherein: the first data structure has a plurality of file management structures; and rebuilding the information related to the first data structure includes using one or more of the plurality of copies of the second data structure to find other file management structures after data corruption of one of the plurality of copies of the second data structure.

7. (Previously Presented) The method as claimed in claim 1 wherein the storage device includes pre-erased recovery blocks; and defining the information related to a first data structure includes writing the plurality of copies of the second data structure to the pre-erased recovery blocks.

8. (Currently Amended) A method for recovering data in a memory comprising:

providing a first base block copy having a first sequence number, a first copy pointer and a first checksum;

providing a second base block copy having a second sequence number, a second copy pointer and a second checksum;

defining a memory location of a data structure within a plurality of base block copies, the plurality of base block copies including the first base block copy and the second base block copy; and

rebuilding the memory location of the data structure by using the first base block copy if the first sequence number is greater than the second sequence number or by using the second base block copy if the second sequence number is greater than the first sequence number.~~using the plurality of base block copies in the event of data corruption of one of the base block copy copies.~~

9. (Currently Amended) The method as claimed in claim 8 wherein the memory performs write operations; and defining a memory location of a data structure within a plurality of base block copies includes updating the first base block copy and the second base block copy ~~plurality of base block copies~~ prior to a write operation.

10. (Cancelled) The method as claimed in claim 8 wherein defining a memory location of a data structure includes differentiating which of the plurality of base block copies has the most recent data.

11. (Previously presented) The method as claimed in claim 8 wherein defining a memory location of a data structure includes selecting each of the plurality of base block copies so at most one of the plurality of base block copies can be corrupted.

12. (Previously presented) The method as claimed in claim 8 wherein: the data structure has an intended state; and rebuilding the location of the data structure includes using one or more of the plurality of base block copies to rebuild the intended state of the data structure after corruption of one of the plurality of base block copies.

13. (Original) The method as claimed in claim 8 wherein the data structure includes pointers to other data structures selected from a group consisting of remap information, wear-leveling tables, configuration data, recovery information, and a combination thereof.

14. (Previously presented) The method as claimed in claim 8 wherein the memory includes pre-erased recovery blocks; and defining a memory location of a data structure includes writing the plurality of base block copies to the pre-erased recovery blocks.

15. (Currently Amended) A method for recovering data in a non-volatile memory comprising:

providing a first base block copy having a first sequence number, a first copy pointer and a first checksum;

providing a second base block copy having a second sequence number, a second copy pointer and a second checksum;

defining in the non-volatile memory a location of a data structure ~~with~~ in at least two base block copies, including the first base block copy and the second base block copy, by storing a first pointer in the first base block copy and a second pointer in the second base block copy, the first pointer and the second pointer each identifying the location of the first data structure, and by setting the first sequence number to represent the number of times the first pointer has been updated successfully in the first base block copy and setting the second sequence number to represent the number of times the second pointer has been has been updated successfully in the second base block copy; and

rebuilding the location of the data structure by using the first base block copy if the first sequence number is greater than the second sequence number or by using the second base block copy if the second sequence number is greater than the first sequence number ~~in the event one base block copy cannot be located or verified by using another base block copy~~.

16. (Original) The method as claimed in claim 15 wherein the non-volatile memory performs write operations; and the base block copies are updated prior to a write operation.

17. (Original) The method as claimed in claim 15 wherein rebuilding the location of the data structure includes differentiating which base block copy has the most recent data.

18. (Original) The method as claimed in claim 15 wherein defining the location of the data structure includes selecting each of the base block copies so at most one can be corrupted.

19. (Previously presented) The method as claimed in claim 15 wherein: the non-volatile memory has an intended state; and at least one of the base block copies can be used to rebuild the intended state of the non-volatile memory before corruption thereof.

20. (Original) The method as claimed in claim 15 wherein the data structure includes pointers to other data structures selected from a group consisting of remap information, wear-leveling tables, configuration data, recovery information, and a combination thereof.

21. (Original) The method as claimed in claim 15 wherein the non-volatile memory includes pre-erased recovery blocks; and the base block copies are written to pre-erased recovery blocks.

22. (Previously presented) The method as claimed in claim 1, wherein defining in the storage device information related to the first data structure within a plurality of copies of a second data structure includes associating a memory address to the first data structure.

23. (Cancelled) The method as claimed in claim 22, wherein defining in the storage device information related to the first data structure within a plurality of copies of a second data structure includes creating a first base block copy and a second base block copy.

24. (New) The method in claim 1, wherein the defining of the storage device information includes storing a first pointer in the first base block copy and a second pointer in the second base block copy, the first pointer and second pointer for identifying the location of the first data structure.

25. (New) The method in claim 24, further including setting the first copy pointer to point to the location of the second base block copy.

26. (New) The method in claim 24, further including: setting the second copy pointer to point to the location of the first base block copy.

27. (New) The method in claim 24, further including: setting the first sequence number to represent the number of times the first base block copy has been updated successfully.

28. (New) The method in claim 24, further including: setting the second sequence number to represent the number of times the second base block copy has been updated successfully.

29. (New) The method in claim 24, further including: setting the first sequence number to represent the number of times the first pointer has been updated successfully.

30. (New) The method in claim 24, further including: setting the first sequence number to represent the number of times the first pointer has been updated successfully.

31. (New) The method in claim 24, further including: providing an original base block copy, the original base block copy having an original checksum which is compared with the first checksum, and if the original checksum and first checksum match, using the first copy pointer to locate the second base block copy.